

Mark Rader  
mrader@us.ibm.com  
IBM zGrowth Unit (ATS)



# IBM DB2® for z/OS: Data Sharing Technical Deep Dive



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## Assumptions



- This is a technical discussion of DB2 for z/OS Data Sharing topics
- The audience should be familiar with DB2 for z/OS Data Sharing concepts, behavior and benefits, based on
  - Experience with a data sharing environment
  - Recent data sharing education
  - DB2 for z/OS publications or Redbooks®
- Flexible capacity
- Scalability
- High availability
- Dynamic workload balancing

## Acronyms



- CF – Coupling Facility LPAR
  - ICF – Integrated CF, aka Internal CF
- CFRM – CF Resource Management, definitions in CFRM policy
- CFCC – CF Control Code
- CF Links – connectivity between CF LPAR and ‘host’ CECs
  - ISC – fiber links, medium to long distance
  - ICB – copper links, very short distance
  - PSIFB - InfiniBand® links, short (12X IB) to long (1X IB) distance
  - IC – internal, microcode links for ICFs
- XCF – Cross-System Coupling Facility – communication between CECs
- XES – Cross-System Extended Services, z/OS component that manages CFs

## Agenda

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- DB2 Data Sharing
  - Configurations
  - Standard CF interaction
  - Performance monitoring
  - Auto Alter
- Workload growth
  - Lock structure
  - GBPs
  - Changes in configuration
    - CF considerations
- What's New in DB2 10 and DB2 11

## DB2 Data Sharing Starting Configuration



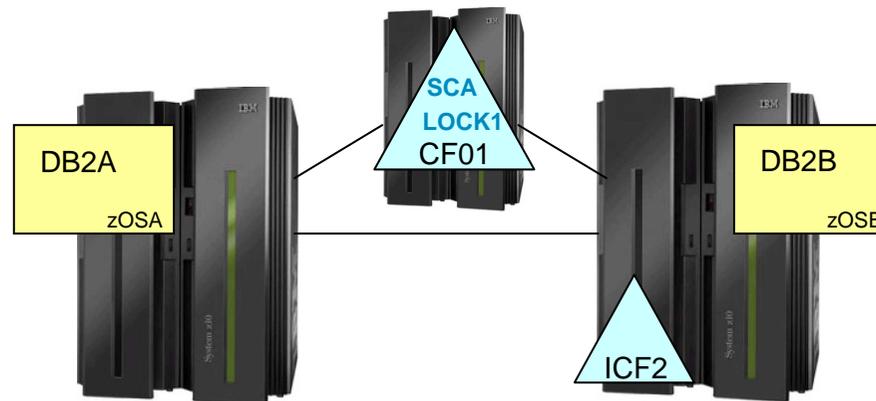
- Starting with DB2 V4



## DB2 Data Sharing: Usual Configuration



- Introduction of ICF

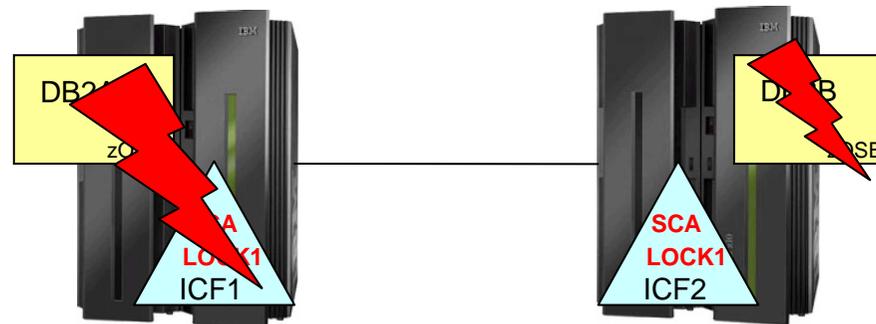


- SCA and LOCK1 on external CF; isolated from DB2 and IRLM members
- Duplexed GBPs spread across CF01 and ICF2

## DB2 Data Sharing: 2-ICF Configuration



- Reduced number of CEC footprints
- Risk of 'double failure': DB2 and SCA, IRLM and LOCK1
  - If structure and exploiter fail, other members fail, too.

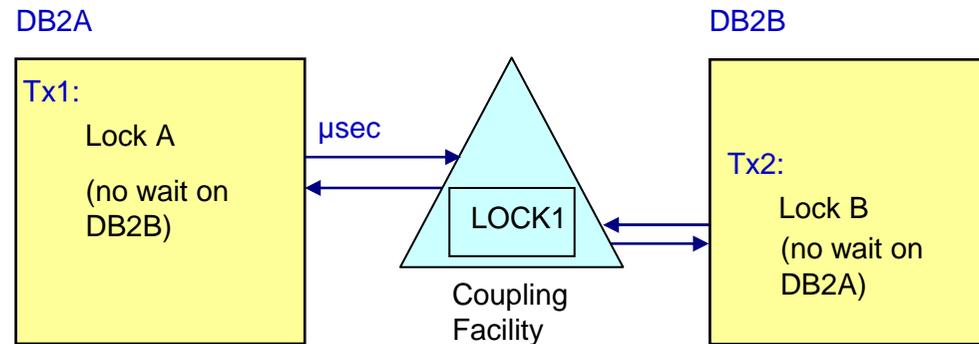


- Duplexed SCA and LOCK1 strongly recommended in this configuration
  - DB2B remains active, even if CEC on left is lost
  - Additional cost: host CPU, CF CPU, and CF link busy 

## Data Sharing: Locking



- Global locking using Parallel Sysplex® coupling technology
  - Inter-system concurrency control



- Cost of obtaining lock does not increase when adding 3<sup>rd</sup> through n<sup>th</sup> members
- This example assumes no contention

## Notes: Lock Structure (LOCK1)

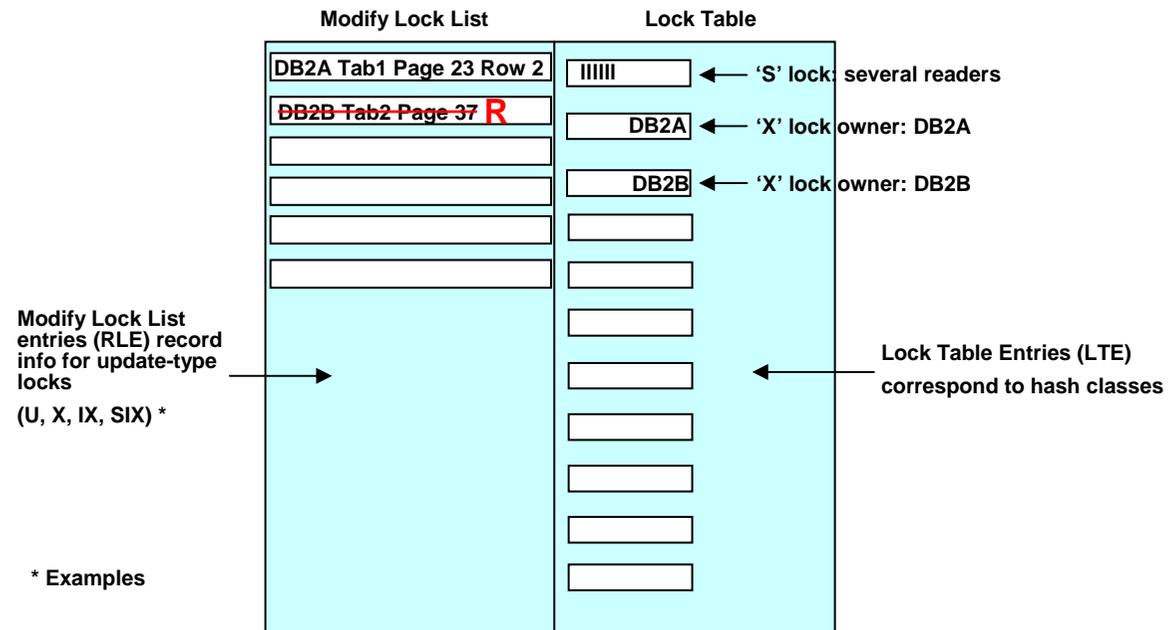


- Used by IRLM to manage global locking
- Holds L-locks and P-locks
  - L-locks to track concurrency
    - Parent L-locks: e.g. table space intent locks
    - Child L-locks: page or row locks
    - Others...
  - P-locks to track coherency. Examples:
    - Page set P-locks: table space, partition, index, index partition
    - Page P-locks: data page (RLL), index leaf page, space map page
    - Others...
- Consists of a lock table (hash table) and a modify lock list
  - Lock table controls access to resources
    - One entry can record multiple readers and one updater (owner)
  - Modify lock list contains detailed information for update-type locks
    - Entries become retained locks in case of an IRLM or DB2 failure

# Lock Structure (LOCK1)



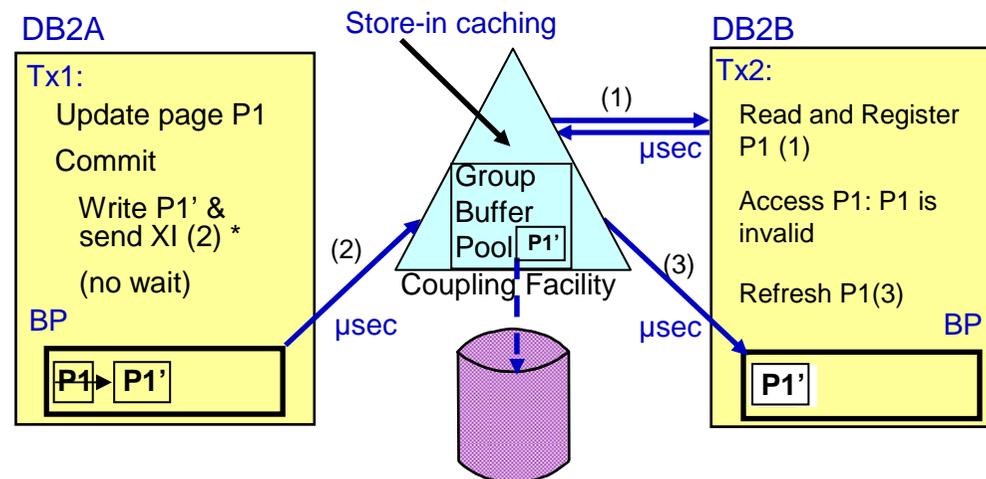
- Simplified view



## Data Sharing: Managing changed data



- Inter-system buffer coherency control
  - Example: DB2A has write interest in the table space, and page P1 is in DB2A's buffer pool



- \* Cross-invalidate (XI) to other member without interrupt

## Notes: Group Buffer Pools (GBPs)



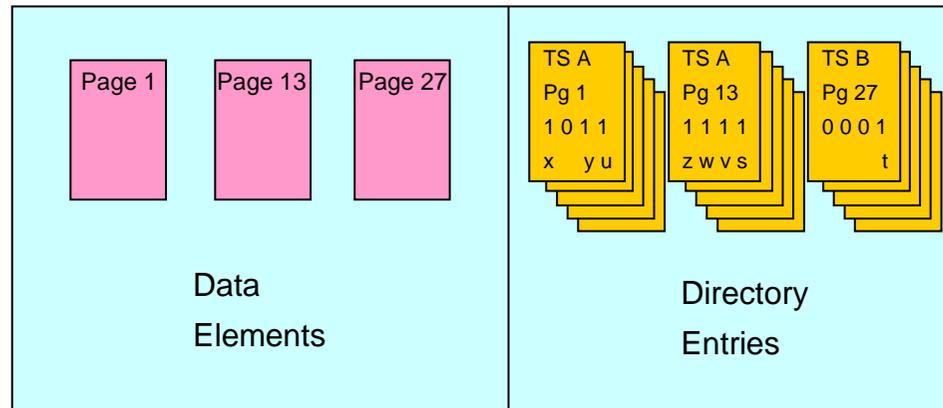
- DB2 uses GBPs to
  - Manage buffer coherency
  - Cache changed pages
    - Optionally cache read-only pages
- GBP consists of directory entries and data elements
  - Directory entries manage coherency by tracking interest in a data or index page by any DB2 member in the data sharing group
    - There is one directory entry for each page in the aggregate pool, no matter how many DB2 members have a copy of that page
  - Data elements are the cached pages that a DB2 member changed
  - In GBP duplexing, data elements exist in both the primary and secondary GBP
    - Directory entries in secondary GBP only exist for the changed pages

## Group Buffer Pool (GBP)



- Simplified view

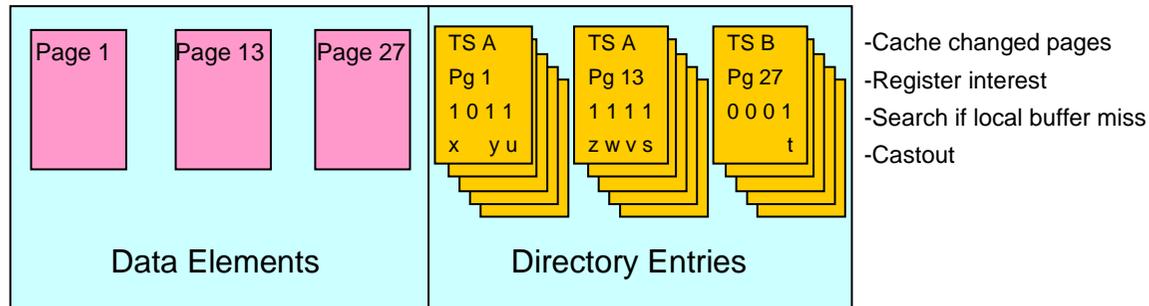
DSNDB20\_GBP2



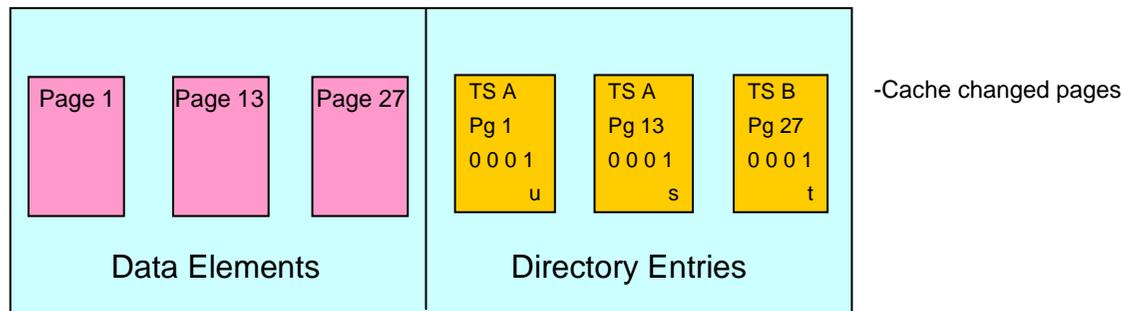
## GBP Duplexing



- DSNDB20\_GBP2 – Primary; “Old” on CF01



- DSNDB20\_GBP2 – Secondary; “New” on ICF2



# Monitoring LOCK1



- RMF CF Activity Report
  - Structure Summary

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COUPLING FACILITY ACTIVITY

PAGE 1

z/OS V1R6                      SYSPLEX \*\*\*\*                      DATE 09/16/2008                      INTERVAL 015.00.489  
 CONVERTED TO z/OS V1R9 RMF                      TIME 08.59.00                      CYCLE 10.000 SECONDS

---

COUPLING FACILITY NAME = CFP01  
 TOTAL SAMPLES(AVG) = 90 (MAX) = 90 (MIN) = 89

---

COUPLING FACILITY USAGE SUMMARY

---

STRUCTURE SUMMARY

TYPE	STRUCTURE NAME	STATUS CHG	★		# REQ	★			LST/DIR ENTRIES	DATA ELEMENTS	LOCK ENTRIES	DIR REC/ XI'S
			ALLOC SIZE	% OF CF STOR		% OF ALL REQ	% OF CF UTIL	AVG REQ/ SEC				
LOCK	DSNDB2B_LOCK1	ACTIVE	16M	0.1	0	0.0	0.0	0.00	24K	0	4194K	N/A
			32						32	0	7381	N/A
	DSNDB2P_LOCK1	ACTIVE	64M	0.5	1483K	10.5	0.0	1646.5	100K	0	17M	N/A
									2121	0	207K	N/A
	DSNDB2Q_LOCK1	ACTIVE	16M	0.1	0	0.0	0.0	0.00	24K	0	4194K	N/A
									272	0	48K	N/A
	DSNDB2R_LOCK1	ACTIVE	16M	0.1	0	0.0	0.0	0.00	24K	0	4194K	N/A
									13	0	5717	N/A

## Notes: Key Points – LOCK1 Structure Summary



- Size can be an issue
  - Determines the number of Lock Table entries (LTE) and space for Modify Lock List entries (RLE)
- Requests per second is important
  - “Busy” is relative; < 5K/sec is not very busy
  - Observed: 166K/sec – very busy
- LIST/DIR ENTRIES = Modify Lock List entries (RLE)
- LOCK ENTRIES = 2-byte Lock Table entries (LTE)
  - May be 4- or 8-byte entries if > 7 members in the data sharing group
  - IRLM automatically rebuilds the lock structure when the 8<sup>th</sup> member (4-byte entries) or 23<sup>rd</sup> member (8-byte entries) joins the data sharing group

## Monitoring LOCK1, cont.



- RMF CF Activity Report
  - Structure Activity

STRUCTURE NAME = DSNDB2P\_LOCK1      TYPE = LOCK      STATUS = ACTIVE

NAME	# REQ		REQUESTS				DELAYED REQUESTS				EXTERNAL REQUEST CONTENTIONS		
	TOTAL	AVG/SEC	# REQ	% OF ALL	-SERV TIME(MIC)-AVG	STD_DEV	REASON	# REQ	% OF REQ	AVG /DEL	REQ TOTAL	REQ DEFERRED	
SYSA	567K		523K	35.3	44.6	64.3	NO SCH	316	0.1	16.8	94.5	0.0	784K
	630.1		44K	3.0	150.0	325.8	PR WT	0	0.0	0.0	0.0	0.0	4634
			0	0.0	INCLUDED	IN ASYNC	PR CMP	3016	0.5	643.6	1418	3.4	4198
													742
SYSB	916K		853K	57.6	43.5	85.3	NO SCH	49	0.0	80.7	184.8	0.0	1256K
	1017		62K	4.2	147.7	259.6	PR WT	0	0.0	0.0	0.0	0.0	5437
			0	0.0	INCLUDED	IN ASYNC	PR CMP	0	0.0	0.0	0.0	0.0	4703
													705
	1483K		1376K	92.8	43.9	78.0	NO SCH	365	0.0	25.3			2040K
	1647		106K	7.2	148.6	288.9	PR WT	0	0.0	0.0			10K
			0	0.0			PR CMP	3016	0.2	933.9	1597	1.9	8901
													1447



## Notes: Key Points – LOCK1 Structure Activity

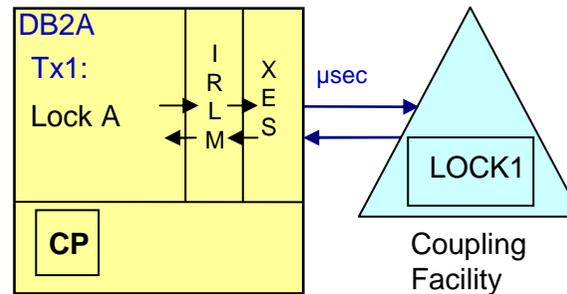


- # REQ TOTAL
  - These are requests on the subchannel
  - Compare with EXTERNAL REQUEST CONTENTIONS: REQ TOTAL, which reflects API requests to XES and should be the higher number
- SERV TIME(MIC) – service time in microseconds
  - ★ – SYNC is key metric – ‘good’ number is relative to CF configuration
    - If ASYNC is non-zero it could be ‘block unlock’, or some requests were converted, either due to subchannel busy or heuristic algorithm
- CONT and FALSE CONT
  - ★ – Contention - recommend:  $\text{CONT}/\text{REQ TOTAL} < 2\%$
  - False Contention - recommend:  $\text{FALSE CONT}/\text{REQ TOTAL} < 1\%$ 
    - If higher, adjust size of LOCK1 to double size of Hash Table

## Heuristic Algorithm and LOCK1



- Most LOCK1 requests are synchronous CF requests
  - Synchronous CF request means host CP is busy for duration of request
  - Long synchronous service times = high host CPU overhead



- XES can convert synchronous request to asynchronous
  - Heuristic algorithm based on measured lock service times
  - Host CP can now do other work during CF request
  - There is some host CP cost to setting up asynchronous request
  - Also elapsed time impact on lock requests

# Monitoring GBPs

- RMF CF Activity Report



- Structure Summary

1 COUPLING FACILITY ACTIVITY PAGE 1

z/OS V1R6 SYSPLEX \*\*\*\* DATE 09/16/2008 INTERVAL 015.00.489  
 CONVERTED TO z/OS V1R9 RMF TIME 08.59.00 CYCLE 10.000 SECONDS

COUPLING FACILITY NAME = CFP01  
 TOTAL SAMPLES(AVG) = 90 (MAX) = 90 (MIN) = 89

COUPLING FACILITY USAGE SUMMARY

STRUCTURE SUMMARY

TYPE	STRUCTURE NAME	STATUS CHG	ALLOC SIZE	% OF CF STOR	# REQ	% OF ALL REQ	% OF CF UTIL	AVG REQ/ SEC	LST/DIR	DATA	LOCK	DIR REC/
									ENTRIES TOT/CUR	ELEMENTS TOT/CUR	ENTRIES TOT/CUR	DIR REC XI'S
CACHE	DSNDB2P_GBP0	ACTIVE	34M	0.3	529	0.0	0.0	0.59	29K	5732	N/A	0
	DSNDB2P_GBP1	ACTIVE	501M	4.1	18380	0.1	0.0	20.41	10	19	N/A	0
	DSNDB2P_GBP16K0	ACTIVE	8M	0.1	120	0.0	0.0	0.13	494K	82K	N/A	0
	DSNDB2P_GBP16K1	ACTIVE	32M	0.3	42641	0.3	0.0	47.35	5406	5406	N/A	0
	DSNDB2P_GBP2	ACTIVE	2G	16.8	8681	0.1	0.0	9.64	1590	1270	N/A	0
	DSNDB2P_GBP3	ACTIVE	8M	0.1	94	0.0	0.0	0.10	0	0	N/A	0
	DSNDB2P_GBP32K	ACTIVE	10M	0.1	132	0.0	0.0	0.15	844	844	N/A	0
	DSNDB2P_GBP32K1	ACTIVE	16M	0.1	120	0.0	0.0	0.13	6008	1201	N/A	0
	DSNDB2P_GBP5	ACTIVE	256M	2.1	358	0.0	0.0	0.40	0	0	N/A	0
									1438	2862	N/A	0
									1	0	N/A	0
									521K	13K	N/A	0
									7	7	N/A	021

## Notes: Key Points – GBPs Structure Summary



- Size and requests per second important
- LIST/DIR ENTRIES = directory entries
- DATA ELEMENTS = data pages
  - If current directory entries = current data pages, probably secondary GBP (GBP duplexing)
    - Could also be the effect of Auto Alter
- DIR REC/DIR REC XI'S = directory reclaims / cross-invalidations (XI's) due to directory reclaims
  - Should be zero! Investigate if non-zero, especially XI's
    - If DIR REC XI'S non-zero, potential performance impact
  - CF report does not have directory reclaim details
    - Use –DIS GBPOOL GDETAIL

## Monitoring GBPs, cont.



- RMF CF Activity Report
  - Structure Activity

```

STRUCTURE NAME = DSNDB2P_GBP1      TYPE = CACHE  STATUS = ACTIVE
-----
SYSTEM # REQ  ----- REQUESTS ----- DELAYED REQUESTS -----
NAME   TOTAL #   % OF -SERV TIME(MIC)- REASON #   % OF --- AVG TIME(MIC) ---
      AVG/SEC REQ ALL  AVG  STD_DEV  # REQ  REQ /DEL  STD_DEV /ALL
-----
SYSP   124  SYNC  19   0.1  12.0   6.8   NO SCH  34   0.3  853.4   837.8   2.3
      ASYNC 12K  67.6  73.2  297.3  PR WT  0   0.0   0.0     0.0   0.0
      CHNGD  0   0.0  INCLUDED IN ASYNC  PR CMP 0   0.0   0.0     0.0   0.0
      DUMP   0   0.0   0.0     0.0   0.0
-----
SYSP   59  SYNC  44   0.2  54.9  246.0  NO SCH  0   0.0   0.0     0.0   0.0
      ASYNC 5887 32.0  54.7  110.9  PR WT  0   0.0   0.0     0.0   0.0
      CHNGD  0   0.0  INCLUDED IN ASYNC  PR CMP 0   0.0   0.0     0.0   0.0
      DUMP   0   0.0   0.0     0.0   0.0
-----
TOTAL 18380 SYNC  63   0.3  41.9  205.9  NO SCH  34   0.2  853.4   837.8   1.6  -- DATA ACCESS --
      20.42 ASYNC 18K  100  67.3  253.0  PR WT  0   0.0
      CHNGD  0   0.0  PR CMP 0   0.0
      DUMP   0   0.0
-----

```

Secondary GBP

## Monitoring GBPs, cont.



- RMF CF Activity Report
  - Structure Activity

```

STRUCTURE NAME = DSNDB2P_GBP1      TYPE = CACHE  STATUS = ACTIVE
# REQ  -----  REQUESTS  -----  DELAYED REQUESTS  -----
SYSTEM  TOTAL  #  % OF  -SERV TIME(MIC)-  REASON  #  % OF  --- AVG  TIME(MIC)  ---
NAME  AVG^SEC  REQ  ALL  AVG  STD_DEV  #  REQ  /DEL  STD_DEV  /ALL
-----
          SYNC  183K  30.5  18.0  47.5  NO SCH  819  0.4  1341  2830  5.6
          ASYNC  14K   2.3  152.7  369.3  PR WT  0  0.0  0.0  0.0  0.0
          CHNGD  818   0.1  INCLUDED IN ASYNC  PR CMP  0  0.0  0.0  0.0  0.0
                                     DUMP  0  0.0  0.0  0.0  0.0
-----
          SYNC  363K  60.5  15.9  74.5  NO SCH  124  0.0  770.1  912.6  0.2
          ASYNC  40K   6.6  143.7  354.7  PR WT  0  0.0  0.0  0.0  0.0
          CHNGD  124   0.0  INCLUDED IN ASYNC  PR CMP  0  0.0  0.0  0.0  0.0
                                     DUMP  0  0.0  0.0  0.0  0.0
-----
TOTAL  599K  SYNC  545K  90.9  16.6  66.7  NO SCH  943  0.2  1266  2665  2.0  -- DATA ACCESS --
          665.9  ASYNC  53K   8.9  146.0  358.5  PR WT  0  0.0
          CHNGD  942   0.2  PR CMP  0  0.0
                                     DUMP  0  0.0
    
```

Primary GBP

## Notes: Key Points – GBPs Structure Activity

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- SERV TIME(MIC)
  - SYNC is key metric – ‘good’ number is relative to CF configuration
    - If REQ/SEC < 100, variations in service time probably not significant
  - ASYNC requests are expected, especially in secondary GBPs
- XI’s in lower right are not necessarily reclaims
  - Most likely business as usual

## Monitoring GBPs: -DIS GBPOOL



- –DIS GBPOOL(\*) TYPE(GCONN) GDETAIL(\*)
  - Contains status and definition information as well as statistics
  - Reports statistics since GBP allocation
- –DIS GBPOOL(\*) TYPE(GCONN) GDETAIL( INTERVAL )
  - To monitor an interval, execute this command before and after the desired interval.
  - Output messages from second command will show GBP statistics for the interval
- Typical problems due to incorrectly defined GBP
  - Directory entry reclaims
  - XIs due to directory entry reclaims
  - Writes failed due to lack of storage

# -DIS GBPOOL(\*) TYPE(GCONN) GDETAIL(\*)



07.57.32	STC34822	DSNB784I	-DB2A GROUP DETAIL STATISTICS	362	
	362		READS		
	362		DATA RETURNED		= 1842830
07.57.32	STC34822	DSNB785I	-DB2A DATA NOT RETURNED	363	
	363		DIRECTORY ENTRY EXISTED		= 1490516
	363		DIRECTORY ENTRY CREATED		= 9995482
	363		DIRECTORY ENTRY NOT CREATED		= 26712646, 0
07.57.32	STC34822	DSNB786I	-DB2A WRITES	364	
	364		CHANGED PAGES		= 50473770
	364		CLEAN PAGES		= 3408467
	364		FAILED DUE TO LACK OF STO		48
	364		CHANGED PAGES SNAPSHOT VALUE		= 5568
07.57.32	STC34822	DSNB787I	-DB2A RECLAIMS	365	
	365		FOR DIRECTORY ENTRIES		80726
	365		FOR DATA ENTRIES		= 28878053
	365		CASTOUTS		= 28679918
07.57.32	STC34822	DSNB788I	-DB2A CROSS INVALIDATIONS	366	
	366		DUE TO DIRECTORY RECLAIMS		56680
	366		DUE TO WRITES		= 2666240
	366		EXPLICIT		= 0
07.57.32	STC34822	DSNB762I	-DB2A DUPLEXING STATISTICS FOR GBP11-SEC	367	
	367		WRITES		
	367		CHANGED PAGES		= 50072797
	367		FAILED DUE TO LACK OF STORAGE		= 48
	367		CHANGED PAGES SNAPSHOT VALUE		= 5568

## Notes: Sizing CF Structures



- <http://www.ibm.com/systems/support/z/cfsizer>
  - CF Structure Sizer Tool
- *DB2 Version 9.1 for z/OS Installation Guide*, GC18-9846
- *DB2 10 for z/OS Installation and Migration Guide*, GC19-2974
- *DB2 11 for z/OS Installation and Migration Guide*, GC19-4056
  - Knowledge Center: [cf sizing for DB2 10 or 11](#)
- Rule of thumb for GBPs
  - Start with CFSizer INITSIZE
  - Round up
  - Make that result INITSIZE; make SIZE up to twice that value
  - Use Auto Alter

## Auto Alter – What is it?



- Autonomic effort by XES to avoid filling up any kind of structure. For GBPs:
  - If all data elements (pages) are changed, writes cannot occur
  - If all directory entries are marked changed, new pages cannot be registered
- Auto Alter has algorithms that
  - can increase or decrease number of entries and/or elements to avoid structure full conditions
  - can increase or decrease the size of the structure
- Can alter, dynamically, the precise directory to data ratio for GBPs
- Design point is for gradual growth, not spikes



## Auto Alter and DB2



- DB2 Structures support Auto Alter
- LOCK1 – effective on Modify Lock List entries (RLEs)
  - Lock Table entries (LTE) cannot be changed without a rebuild
- SCA – can be increased
- Main value is for Group Buffer Pools (GBPs). Why?
  - People tend not to tune GBPs
    - Organizational division of labor
      - DB2 DBAs responsible for local BPs – may forget about GBPs
      - z/OS responsible for GBPs – and they own the CFRM Policy
    - DB2 needs ?? more directory entries than data page elements
    - Each –ALTER to change directory entries means manual GBP rebuild
- Works for duplexed GBPs

## Auto Alter – When not to use it



- CF available storage is <10%
  - Auto Alter reduces the size of “alterable” structures below INITSIZE (to MINSIZE), attempting to get 10% available storage in the CF
- Not enough storage for size of structure, especially in Test environments
  - XES reaches SIZE quickly
  - Reclaim avoidance results in constant XES attempts to increase directory entries and reduce data pages
    - Reclaim avoidance alone does not allow structure size increase
  - Attempts usually fruitless - produce alarming console messages
  - Hint: test one structure, correctly sized, instead of all

## Workload Growth

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- Increased transaction, batch and/or query volumes
- New applications
- Mergers
- New business opportunities
- Regulatory compliance
- Technology advances

## Workload Changes and LOCK1



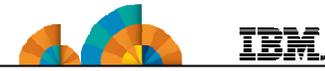
- Increased lock requests may lead to
  - Higher CF CPU busy
  - Higher synchronous service time, and host CPU cost
  - Higher transaction or query elapsed time, higher job run time
- New applications may not follow standards
  - Less lock avoidance by new applications may mean more locking for existing applications
  - Long commit scopes hold Modify Lock List entries (RLEs) longer
  - Row level locking increases demand for RLEs
- False contention could increase, requiring more Lock Table Entries (LTEs)

## Workload Changes and LOCK1



- Possible solutions
  - Increase CF CPU capacity
    - More CPs and/or faster CPs
  - Increase the number of RLEs
    - SETXCF START,ALTER,strnm=&,SIZE=&  
to increase the size of LOCK1
      - Assumes allocation < SIZE in CFRM policy
      - Else change CFRM policy definition, rebuild structure
  - Increase the number of LTEs
    - Requires a structure rebuild with larger allocation
      - CFRM policy change required if allocation already = SIZE
  - CF storage increase may be necessary

# LOCK1 Example



- RMF CF Activity Report
  - Structure Activity

STRUCTURE NAME = DSN\*\*\*\*\_LOCK1    TYPE = LOCK    STATUS = ACTIVE

SYSTEM NAME	# REQ		REQUESTS				REASON	DELAYED REQUESTS					EXTERNAL REQUEST	
	TOTAL	AVG/SEC	# REQ	% OF ALL	-SERV TIME(MIC)-AVG	STD_DEV		# REQ	% OF REQ	/DEL	AVG TIME(MIC)	STD_DEV	/ALL	REQ TOTAL
S***	232M		232M	38.8	11.0	5.8	NO SCH	88K	0.0	30.8	236.2	0.0	REQ TOTAL	133M
	64403		ASYNC	750	0.0	413.5	PR WT	0	0.0	0.0	0.0	0.0	REQ DEFERRED	239K
			CHNGD	0	0.0		PR CMP	0	0.0	0.0	0.0	0.0	-CONT	239K
													-FALSE CONT	162K
S***	187M		187M	31.2		5.5	NO SCH	14K	0.0	15.0	67.2	0.0	REQ TOTAL	99M
	51870		ASYNC	0	0.0	0.0	PR WT	0	0.0	0.0	0.0	0.0	REQ DEFERRED	218K
			CHNGD	0	0.0		PR CMP	0	0.0	0.0	0.0	0.0	-CONT	218K
													-FALSE CONT	111K
S***	179M		179M	30.0		5.9	NO SCH	50	0.0	19.7	37.5	0.0	REQ TOTAL	106M
	49841		ASYNC	1500	0.0	87.7	PR WT	0	0.0	0.0	0.0	0.0	REQ DEFERRED	333K
			CHNGD	0	0.0		PR CMP	0	0.0	0.0	0.0	0.0	-CONT	333K
													-FALSE CONT	242K
	598M		598M	100	10.7	5.7	NO SCH	108K	0.0	28.2	215.6	0.0	REQ TOTAL	337M
	166.1K		ASYNC	2250	0.0	80.2	PR WT	0	0.0	0.0	0.0	0.0	REQ DEFERRED	790K
			CHNGD	0	0.0		PR CMP	0	0.0	0.0	0.0	0.0	-CONT	790K
													-FALSE CONT	515K

Annotations: A blue star is placed below the first row of the last system. A large blue arrow points down from the '11.0' value in the first system's '% OF ALL' column. A large pink arrow points down from the '14K' value in the second system's '# REQ' column. A large orange arrow points right from the '598M' value in the third system's 'TOTAL' column.

## Workload Changes and GBPs

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- Increased GBP requests may lead to
  - Higher CF CPU busy
  - Higher synchronous service time, and host CPU cost
  - Higher transaction or query elapsed time, higher job run time
- New applications may
  - Change access patterns of existing tables or indexes
  - Add tables and indexes to existing buffer pools
- Local buffer pool allocations may increase
  - GBPs might be forgotten

## Workload Changes and GBPs

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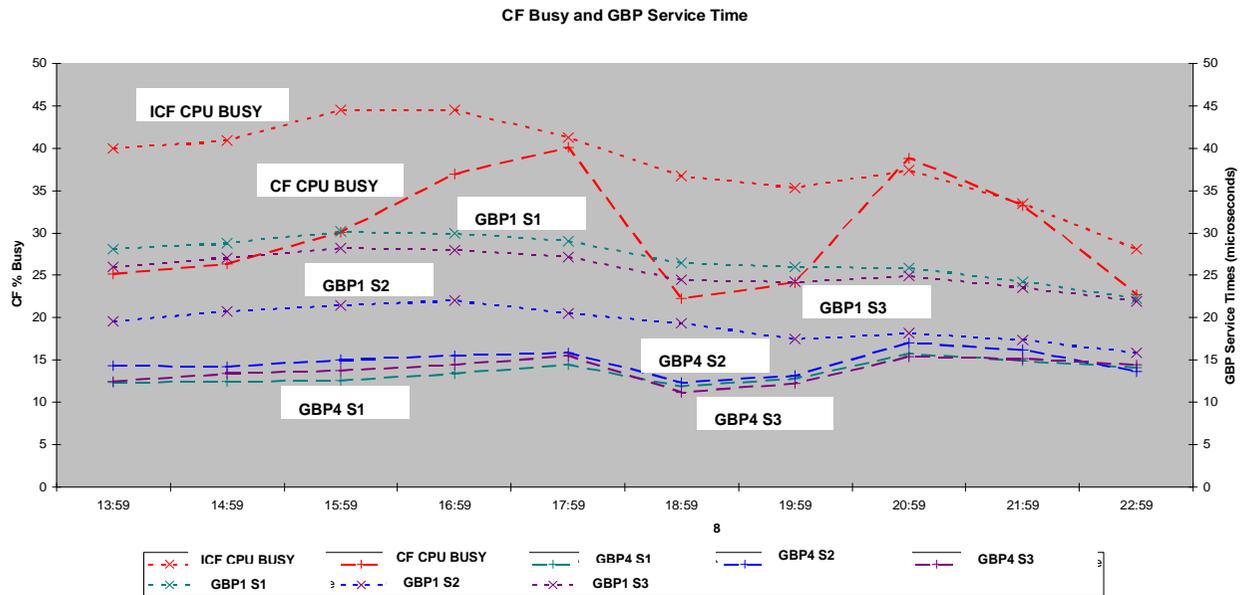


- Possible solutions
  - Increase CF CPU capacity
    - More CPs and/or faster CPs
  - Increase the size of the GBPs
  - Tune local buffer pool thresholds and GBP thresholds
  - CF storage increase may be necessary

# GBPs and Impact of CF Busy



- ICF has two CPs
- CF (external) has three CPs



## GBPs and Impact of CF Busy, cont.



- GBP1 on ICF was very busy over a 10 hour interval
  - S1: 178 M synchronous requests
  - S3: 144.5 M synchronous requests
  - If 10  $\mu$ sec saved from each request, over 300 CPU seconds per hour of 'host effect' could be saved from GBP1 alone
- How could 10  $\mu$ sec be saved?
  - Increase number of CPs on ICF to reduce CF busy and improve service time
  - Upgrade CEC with ICF to reduce CF service times

## Workload Changes and SCA

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- New applications or new workloads may add tables and indexes
- New clients may require additional databases
- Auto Alter may be able to handle most of the increase
- Use CF Structure Sizer Tool to validate CFRM policy definition

## When New Members Join the Data Sharing Group

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- GBP – review sizes
  - Increased demand for directory entries and data elements
  - Auto Alter may not be sufficient to handle multiple new members
- LOCK1
  - 4-byte LTEs required when 8<sup>th</sup> member joins the group
    - Automatic rebuild will normally result in half as many LTEs, so false contention will increase
    - Prepare for larger LTEs before adding 8<sup>th</sup> member
  - 8-byte LTEs required when 23<sup>rd</sup> member joins the group
    - Automatic rebuild has same considerations

## Configuration Changes

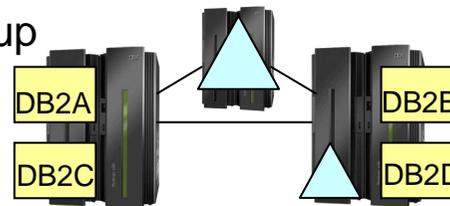


- CF Considerations
  - Balanced performance: CF technology = CEC technology
  - Unbalanced configuration examples:
    - zEC12 CF and z196 CEC – good for the CEC
    - z196 CF and zEC12 CEC – more Host Effect cost to CEC
    - z10 CF and zEC12 CEC – ‘heuristic algorithm’ likely to convert many synchronous requests to asynchronous
      - Algorithm represents tradeoff of host effect versus cost of conversion
      - Elapsed times, contention, and time outs likely to increase
  - Increase in distance between CF and CEC can have similar effect
    - Asynchronous conversion frequently observed as distance between CEC and CF increases

## Creative Use of CF Storage



- As more DB2 members join the group
  - Consider GBPCACHE ALL
    - Each page is read into GBP on first access
    - Only one member incurs I/O cost for each page
    - Local buffers can be smaller – GBP acts as very fast cache
- If large objects with very random access and minimal page re-reference
  - Consider GBPCACHE NONE
    - Saves GBP access on local page miss
    - Enforces ‘store through cache’: synchronous writes to disk at commit
      - Modern cache controllers minimize negative impact





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## **DB2 10 for z/OS**

### **Data Sharing Highlights**

## DB2 10 for z/OS and Data Sharing



- Deleting member of data sharing group
    - Offline utility
  - Deleting structures during group restart
    - DEL\_CFSTRUCTS\_ON\_RESTART - DSNZPARM for DR
  - Sub-group attach
  - DDF Restart Light – handle indoubts
  - MEMBER CLUSTER for UTS
  - -MODIFY DDF – online changes for LOCATION ALIAS
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- LRSN spin avoidance
  - IFCID 359 – index split
  - GBP DELETE\_NAME processing
  - BP scan avoidance



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# DB2 11 for z/OS

## Data Sharing Highlights

## DB2 11 for z/OS Data Sharing Enhancements



- Castout enhancements: New CLASST setting – similar to VDWQT
- RESTART LIGHT Enhancements
- Buffer pool enhancements
  
- GBP Write-around
  - If GBP / CF busy, write new pages to directly to disk
  - Reduce impact of flood of new pages on rest of GBP
- Automatic LPL or GRECP recovery
- CF DELETE\_NAME
- Locking enhancements
- Index split performance
- LRSN spin avoidance – extended LRSN

## Additional Resources

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- ***Data Sharing: Planning and Administration***
  - DB2 9 for z/OS: SC18-9845
  - DB2 10 for z/OS: SC19-2973
  - DB2 11 for z/OS: SC19-4055
  - [KC db2 data sharing planning](#)



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Questions?



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**Thank you!**